

Dynamic equilibrium framework in cyber-security game

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Outline

- ▶ Goal & Motivation
- ▶ Game model
- ▶ Literature
- ▶ Characterization of equilibrium
- ▶ Results

Motivation

- ▶ Number of cyber threats is increasing
 - ▶ annual cost caused by cyber crimes increased by 19% in USA (2015)
 - ▶ 99% of companies experienced cyber attack in 2015
- ▶ Needs for cyber-security related insurance is growing
 - ▶ \$2 billion in 2014 → \$5 billion by 2018
- ▶ But cyber-insurance market is premature
 - ▶ insurance rates are fluctuating
 - ▶ hard to measure cyber risks
 - ▶ **strategic cyber attacker**

Goal

- ▶ Construct a game model:
 - ▶ strategic attacker and defender
 - ▶ financially motivated players
 - ▶ containing some specific features of cyber-security
 - ▶ general enough (no CS-technology)
 - ▶ tractable
- ▶ Analyze the game:
 - ▶ existence, uniqueness of equilibrium
 - ▶ equilibrium with strategic attacker, non-strategic attacker
 - ▶ roll of defender

Game model: roll of the players

- ▶ **Defender** (example: Internet service provider)
 - ▶ can observe user's action with noise
 - ▶ doesn't know whether the user is an attacker or not
 - ▶ can update the probability that the user is an attacker
 - ▶ **can block the user** based on the suspicion level
 - ▶ chooses the **optimal threshold** to minimize the expected costs
- ▶ **User** (example: IP address)
 - ▶ can be an attacker or a normal user
 - ▶ normal user: no malicious activity
 - ▶ attacker: malicious activity → profit
 - ▶ too much malicious activity → blocked by defender **sooner**
 - ▶ attacker chooses optimal activity level to maximize the expected profit

Game model: parameters & notation

- ▶ $\theta \in \{0, 1\}$: identity of user ($\theta = 0$ for normal user, $\theta = 1$ for attacker)
- ▶ $q_0 = \mathbb{E}[\theta] \in (0, 1)$: Initial suspicion level
- ▶ If attacker's action is given by stream $\alpha_t dt$, the defender can observe Y_t with

$$dY_t = \alpha_t 1_{\{\theta=1\}} dt + dW_t$$

where W_t is a standard Brownian motion, independent of θ .

- ▶ Defender continuously updates suspicion level (q .) based on Y .

$$q_t = \mathbb{E}[\theta | \mathcal{F}_t^Y]$$

- ▶ If defender chooses blocking threshold $p \in [0, 1]$, the game is over at

$$\tau = \inf \{t \geq 0 : q_t \geq p\}$$

- ▶ Attacker believes that the q . process obeys following SDE:

$$dq_t = \lambda(q_t) (dY_t - \mu(q_t) dt)$$

Game model: equilibrium

Definition of equilibrium: $(\lambda, \mu, \alpha, \hat{p})$ is called an equilibrium if following holds. $(\lambda, \mu, \alpha : \mathbb{R}^2 \rightarrow \mathbb{R}, \hat{p} \in [0, 1].)$

(1) (Attacker's optimization) Let $p \in [0, 1]$ be given. Assume that attacker **believes** that when his attacking intensity is Δ_t , then the q_t obeys the following SDE:

$$dq_t = \lambda(p, q_t)(\Delta_t dt + dW_t - \mu(p, q_t)dt).$$

Then, α solves attacker's maximization problem:

$$\alpha(p, q.) = \arg \max_{0 \leq \Delta_t \leq M} \mathbb{E} \left[\int_0^\tau e^{-rt} \Delta_t dt \right]$$

$$\text{where } \tau := \inf\{t \geq 0 : q_t \geq p\}$$

r : discounting factor

M : limit of attacker's action

Game model: equilibrium

(2) (Defender's optimization) Let q_t be the solution of following SDE:

$$dq_t = \lambda(p, q_t)(dY_t - \mu(p, q_t)dt)$$

where $dY_t = \alpha(p, q_t)1_{\{\theta=1\}}dt + dW_t$

Then q_t agrees with suspicion level, i.e., $q_t = \mathbb{E}[\theta | \mathcal{F}_t^Y]$ holds.
And \hat{p} solves the defender's cost-minimization problem:

$$\hat{p} = \arg \max_{p \in [0,1]} \mathbb{E} \left[\int_0^\tau (e^{-rt} \alpha(p, q_t) \cdot 1_{\{\theta=1\}}) dt + e^{-r\tau} l \cdot 1_{\{\theta=0\}} \right]$$

where $\tau := \inf\{t \geq 0 : q_t \geq p\}$

where $l > 0$ is an one-time-cost of blocking an innocent user.

Related literature

- ▶ Insider trading models
 - ▶ game between insider(private info) and market maker
 - ▶ Kyle (1985), Back (1992), etc.
 - ▶ **our model**: market maker can stop the game
- ▶ Attacker-Defender framework
 - ▶ game between attacker and **myopic** defender
 - ▶
 - ▶ **our model**: defender is not myopic and stops the game
- ▶ Sequential hypothesis testing
 - ▶ obtain sequence of test results for hypothesis testing
 - ▶ optimal stopping problem to minimize expected costs
 - ▶ Wald (1945), etc.
 - ▶ **our model**: game between test-result-provider and tester

Remarks on our game model

- ▶ **Unique feature**

- ▶ Internet security provider (defender) has authority to block users (ability to stop the game)
- ▶ Defender is not myopic (focused on the defender's roll)

- ▶ **General feature**

- ▶ Attacker's and defender's strategies are driven by financial motives
- ▶ No specific hacking-technology: our model can be used for any warfare-type situation

Characterization of equilibrium

(1) Attacker's optimization problem:

$$V(q_0) := \max_{0 \leq \Delta \leq M} \mathbb{E} \left[\int_0^{\tau} e^{-rt} \Delta_t dt \right]$$

where

$$\begin{cases} dq_t = \lambda(\Delta_t dt + dW_t - \mu dt) \\ \tau := \inf\{t \geq 0 : q_t \geq p\} \end{cases}$$

DPP implies HJB equation:

$$\begin{cases} \max_{\Delta \in [0, M]} \left(-rV(q) - V'(q)\lambda(q)\mu(q) + \frac{1}{2}V''(q)\lambda^2(q) + \Delta(1 + V'(q)\lambda(q)) \right) = 0 \\ V(p) = 0 \\ V(0+) = \frac{M}{r} \end{cases}$$

Characterization of equilibrium

(2) Defender's estimation: The solution of the SDE

$$dq_t = \lambda(q_t)(dY_t - \mu(q_t)dt)$$

where $dY_t = \alpha(q_t)1_{\{\theta=1\}}dt + dW_t$

should satisfy $q_t = \mathbb{E}[\theta | \mathcal{F}_t^Y]$. Filtering equation produces

$$\begin{aligned} dq_t &= \left(\mathbb{E}[\theta \alpha(q_t) 1_{\{\theta=1\}} | Y_{[0,t]}] - \mathbb{E}[\theta | Y_{[0,t]}] \mathbb{E}[\alpha(q_t) 1_{\{\theta=1\}} | Y_{[0,t]}] \right) \\ &\quad \left(dY_t - \mathbb{E}[\alpha(q_t) 1_{\{\theta=1\}} | Y_{[0,t]}] dt \right) \quad (1) \\ &= q_t(1 - q_t)\alpha(q_t)(dY_t - q_t\alpha(q_t)dt) \end{aligned}$$

By comparing SDEs for q_t , we obtain relations among λ, μ, α :

$$\alpha(q) = \frac{\lambda(q)}{(1-q)q}, \quad \mu(q) = \frac{\lambda(q)}{1-q}$$

Characterization of equilibrium

Combine these relations with the attacker's HJB,

$$\begin{cases} \frac{V''(q)}{2} - \frac{V'(q)}{1-q} - rV(q)V'(q)^2 = 0, & \text{if } V'(q) < -\frac{1}{M(1-q)q} \\ \frac{V''(q)}{2} + \frac{V'(q)}{q} - \frac{rV(q)}{M^2q^2(1-q)^2} + \frac{1}{Mq^2(1-q)^2} = 0, & \text{if } V'(q) \geq -\frac{\sigma^2}{M(1-q)q} \\ V(0+) = \frac{M}{r}, \quad V(p) = 0 \end{cases}$$
$$\mu(q) = \frac{\lambda(q)}{1-q}, \quad \alpha(q) = \frac{\lambda(q)}{(1-q)q}$$

Finally, the defender's optimization problem is

$$\mathbb{E}\left[\int_0^{\tau} \left(e^{-rt} \alpha(q_t) \cdot \mathbf{1}_{\{\theta=1\}}\right) dt + e^{-r\tau} l \cdot \mathbf{1}_{\{\theta=0\}}\right] \quad \text{with } \tau := \inf\{t \geq 0 : q_t \geq p\}$$
$$= q_0 V(q_0) + (1 - q_0) u(q_0)$$

where $u(q_0) := \mathbb{E}[e^{-r\tau} | \theta = 0]$. Then u should satisfy

$$\frac{1}{2} u''(q) - \frac{u'(q)}{1-q} - \frac{ru(q)}{\lambda(q)^2} = 0, \quad u(p) = 1, \quad u(0) = 0$$

Results

Let φ, a, b, c, x^* be defined as

$$\varphi(x) := \frac{2}{\sqrt{\pi}} \int_0^x e^{-t^2} dt$$

$$a := \frac{1}{2} \left(\sqrt{1 + \frac{8r}{M^2}} - 1 \right)$$

$$b := \frac{M}{\sqrt{r}} - \frac{\sqrt{r}}{aM}$$

$$c := \frac{2\sqrt{r}e^{-b^2} + M\sqrt{\pi}\varphi(b)}{Mp\sqrt{\pi}}$$

$$x^* := \frac{pc - \varphi(b)}{c - \varphi(b)}$$

Let V be defined as

If $\frac{r}{M^2} \geq 1$:

$$V(x) := \frac{M}{r} \left(1 - \left(\frac{1-p}{p} \right)^a \left(\frac{x}{1-x} \right)^a \right), \quad x \in [0, p]$$

If $\frac{r}{M^2} < 1$:

$$V(x) := \begin{cases} \frac{M}{r} - \frac{\sigma^2}{aM} \left(\frac{(1-p)c}{pc - \varphi(b)} \right)^a \cdot \left(\frac{x}{1-x} \right)^a, & \text{for } 0 \leq x \leq x^* \\ \frac{\sigma}{\sqrt{r}} \varphi^{-1} \left(c \cdot \frac{p-x}{1-x} \right), & \text{for } x^* < x \leq p \end{cases}$$

Results

Theorem

We have the unique equilibrium:

$$\lambda(p, x) = \begin{cases} M(1-x)x, & \text{for } 0 \leq x \leq x^* \\ -\frac{1}{V'(x)}, & \text{for } x^* < x \leq p' \end{cases}$$

$$\mu(p, x) = \frac{\lambda(x)}{1-x}, \quad \alpha(p, x) = \frac{\lambda(x)}{(1-x)x}$$

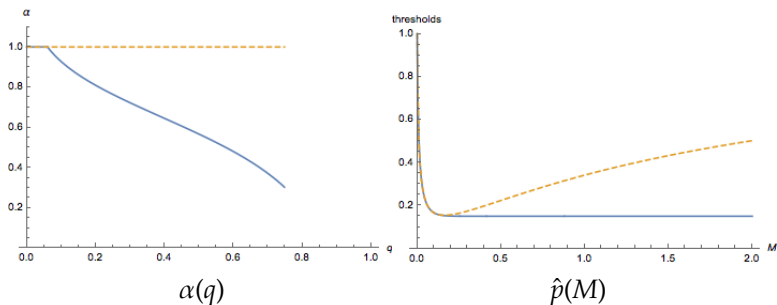
$$\text{If } \frac{r}{M^2} \geq 1 : \quad \hat{p} = \frac{(1+a)rl}{(1+a)rl+aM}$$

$$\text{If } \frac{r}{M^2} < 1 : \quad \hat{p} = \frac{l\sqrt{r}\left(2\sqrt{re^{-b^2}}+M\sqrt{\pi}\varphi(b)\right)}{M+1\sqrt{r}\left(2\sqrt{re^{-b^2}}+M\sqrt{\pi}\varphi(b)\right)}$$

Ingredients of proof

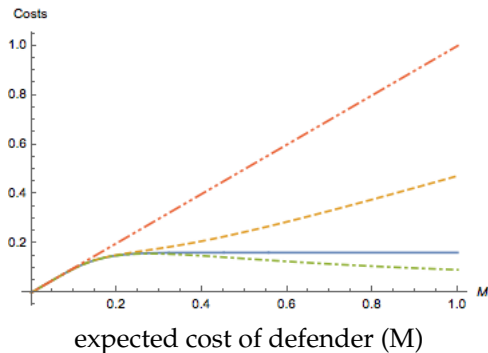
- ▶ Need to check the solutions of SDEs are well defined ($q_t = 0$ is an absorbing state)
- ▶ Check that q cannot reach 0, using Feller's explosion test
- ▶ Verification of optimization problems, filtering equation, stochastic representation $\mathbb{E}[e^{-rt} | \theta = 0]$

Equilibrium strategies



(—): strategic attacker
(- - -): non-strategic attacker

Equilibrium costs comparison with benchmarks



(- · - · - · -): no defender (no blocking at all)

(- · - · -): non-strategic attacker VS defender

(- - -): strategic attacker VS naive defender

(—): strategic attacker VS defender (our model)

Conclusion

- ▶ Dynamic game model for cyber-security
- ▶ Unique feature of non-myopic defender who can optimally stop the game
- ▶ Strategic nature of attacker is important factor
- ▶ Applicable to cyber-insurance market